// HALBORN

Stader Labs -Hedera Stader Protocol v3 Smart Contract Security Audit

Prepared by: Halborn Date of Engagement: Oct 21st, 2022 - Nov 11th, 2022 Visit: Halborn.com

DOCUMENT REVISION HISTORY		
CONTACTS	6	
1 EXECUTIVE OVERVIEW	8	
1.1 INTRODUCTION	9	
1.2 AUDIT SUMMARY	9	
1.3 TEST APPROACH & METHODOLOGY	9	
RISK METHODOLOGY	10	
1.4 SCOPE	12	
2 ASSESSMENT SUMMARY & FINDINGS OVERVIEW	13	
3 FINDINGS & TECH DETAILS	14	
3.1 (HAL-01) TIMELOCK CAN BE BYPASSED - MEDIUM	16	
Description	16	
Code Location	16	
Risk Level	17	
Proof of Concept	18	
Recommendation	18	
Remediation Plan	19	
3.2 (HAL-02) LACK OF TRANSFEROWNERSHIP PATTERN - LOW	20	
Description	20	
Risk Level	20	
Recommendation	20	
Remediation Plan	21	
3.3 (HAL-03) DIFFERENT PROXYNODE ARRAY LENGTH REQUIREME	NTS - LOW	
Description	22	

	Risk Level	24
	Recommendation	24
	Remediation Plan	24
3.4	(HAL-04) MISSING ZERO ADDRESS CHECKS - LOW	25
	Description	25
	Code Location	25
	Risk Level	25
	Recommendation	26
	Remediation Plan	26
3.5	(HAL-05) LACK OF PARAMETER LIMITS - LOW	27
	Description	27
	Code Location	27
	Risk Level	28
	Recommendation	28
	Remediation Plan	29
3.6	(HAL-06) MISSING REENTRANCY GUARD - LOW	30
	Description	30
	Code Location	30
	Risk Level	31
	Recommendation	31
	Remediation Plan	31
3.7	(HAL-07) NODEPROXY ARRAY CANNOT BE MODIFIED - INFORMATIONAL	32
	Description	32
	Code Location	32
	Risk Level	33
	Recommendation	33

	Remediation Plan	33
3.8	(HAL-08) FLOATING PRAGMA - INFORMATIONAL	34
	Description	34
	Risk Level	34
	Recommendation	34
	Remediation Plan	34
3.9	(HAL-09) CACHE ARRAY LENGTH IN FOR LOOPS CAN SAVE GAS - INFO MATIONAL)R- 35
	Description	35
	Code Location	35
	Risk Level	36
	Recommendation	36
	Remediation Plan	36
3.10	(HAL-10) USE CUSTOM ERRORS INSTEAD OF REVERT STRINGS - INFORM	1A- 37
	Description	37
	Risk Level	37
	Recommendation	37
	Remediation Plan	37
3.11	(HAL-11) REVERT STRING SIZE OPTIMIZATION - INFORMATIONAL	38
	Description	38
	Code Location	38
	Risk Level	38
	Recommendation	38
	Remediation Plan	38
3.12	(HAL-12) UNUSED EVENTS - INFORMATIONAL	39
	Description	39

Code Location	39
Risk Level	39
Recommendation	40
Remediation Plan	40
3.13 (HAL-13) UNNECESSARY CHECK - INFORMATIONAL	41
Description	41
Code Location	41
Risk Level	42
Recommendation	42
Remediation Plan	42
3.14 (HAL-14) NO NEED TO INITIALIZE VARIABLES WITH DEFAULT VALU	JES - 43
Description	43
Code Location	43
Risk Level	44
Recommendation	44
Remediation Plan	44
3.15 (HAL-15) USING POSTFIX OPERATORS IN LOOPS - INFORMATIONAL	45
Description	45
Code Location	45
Risk Level	45
Recommendation	45
Remediation Plan	46
3.16 (HAL-16) DIVISION BY ZERO - INFORMATIONAL	47
Description	47
Code Location	47

	Risk Level	47
	Recommendation	47
	Remediation Plan	48
3.17	(HAL-17) SPLITTING REQUIRE() STATEMENTS THAT USES AND OPERATIONAL	TOR 49
	Description	49
	Code Location	49
	Proof of Concept	49
	Risk Level	50
	Recommendation	50
	Remediation Plan	50
4	MANUAL TESTING	51
4.1	INTRODUCTION	52
4.2	TESTING	53
	REWARD MECHANISM	53
	WITHDRAW FUNDS MECHANISM	56
	UNDELEGATION MECHANISM	60
5	AUTOMATED TESTING	64
5.1	STATIC ANALYSIS REPORT	65
	Description	65
	Slither results	65

DOCUMENT REVISION HISTORY

VERSION	MODIFICATION	DATE	AUTHOR
0.1	Document Creation	10/26/2022	Kaan Caglan
0.2	Draft Updates	10/27/2022	Kubilay Onur Gungor
0.3	Draft Updates	10/30/2022	Kaan Caglan
0.4	Draft Updates	11/01/2022	Francisco González
0.5	Draft Review	11/03/2022	Gabi Urrutia
1.0	Remediation Plan	11/04/2022	Francisco González
1.1	Remediation Plan Updates	11/05/2022	Francisco González
1.2	Remediation Plan Review	11/05/2022	Gabi Urrutia

CONTACTS

CONTACT	COMPANY	EMAIL	
Rob Behnke	Halborn	Rob.Behnke@halborn.com	
Steven Walbroehl	Halborn	Steven.Walbroehl@halborn.com	

Gabi Urrutia	Halborn	Gabi.Urrutia@halborn.com	
Kubilay Onur Gungor	Halborn	Kubilay.Gungor@halborn.com	
Kaan Caglan Halborn		Kaan.Caglan@halborn.com	
Francisco González Halborn		Francisco.Villarejo@halborn.com	

EXECUTIVE OVERVIEW

1.1 INTRODUCTION

Stader Labs engaged Halborn to conduct a security audit on their smart contracts beginning on October 21st, 2022 and ending on November 11th, 2022. The security assessment was scoped to the smart contracts provided in the GitHub repositories stader-labs/hedera-stader-protocol-v1

1.2 AUDIT SUMMARY

The team at Halborn was provided a week for the engagement and assigned a full-time security engineer to audit the security of the smart contracts. The security engineer is a blockchain and smart-contract security expert with advanced penetration testing, smart-contract hacking, and deep knowledge of multiple blockchain protocols.

The purpose of this audit is to:

- Ensure that smart contract functions operate as intended.
- Identify potential security issues with the smart contracts.

In summary, Halborn identified some security risks were mostly addressed by the Stader Labs team.

1.3 TEST APPROACH & METHODOLOGY

Halborn performed a combination of manual and automated security testing to balance efficiency, timeliness, practicality, and accuracy in regard to the scope of this audit. While manual testing is recommended to uncover flaws in logic, process, and implementation; automated testing techniques help enhance coverage of the code and can quickly identify items that do not follow the security best practices. The following phases and associated tools were used during the audit:

- Research into architecture and purpose.
- Smart contract manual code review and walkthrough.
- Graphing out functionality and contract logic/connectivity/functions. (solgraph)
- Manual assessment of use and safety for the critical Solidity variables and functions in scope to identify any arithmetic related vulnerability classes.
- Manual testing by custom scripts.
- Scanning of solidity files for vulnerabilities, security hot-spots or bugs. (MythX)
- Static Analysis of security for scoped contract, and imported functions. (Slither)
- Testnet deployment. (Brownie, Remix IDE)

RISK METHODOLOGY:

Vulnerabilities or issues observed by Halborn are ranked based on the risk assessment methodology by measuring the **LIKELIHOOD** of a security incident and the **IMPACT** should an incident occur. This framework works for communicating the characteristics and impacts of technology vulnerabilities. The quantitative model ensures repeatable and accurate measurement while enabling users to see the underlying vulnerability characteristics that were used to generate the Risk scores. For every vulnerability, a risk level will be calculated on a scale of 5 to 1 with 5 being the highest likelihood or impact.

RISK SCALE - LIKELIHOOD

- 5 Almost certain an incident will occur.
- 4 High probability of an incident occurring.
- 3 Potential of a security incident in the long term.
- 2 Low probability of an incident occurring.
- 1 Very unlikely issue will cause an incident.

RISK SCALE - IMPACT

- 5 May cause devastating and unrecoverable impact or loss.
- 4 May cause a significant level of impact or loss.

- 3 May cause a partial impact or loss to many.
- 2 May cause temporary impact or loss.
- 1 May cause minimal or un-noticeable impact.

The risk level is then calculated using a sum of these two values, creating a value of 10 to 1 with 10 being the highest level of security risk.

CRITICAL	HIGH	MEDIUM	LOW	INFORMATIONAL
10 - CRITICAL 9 - 8 - HIGH				
7 - 6 - MEDIUM				
5 - 4 - LOW				
3 - 1 - VERY LO	OW AND INFORMA	TIONAL		

1.4 SCOPE

IN-SCOPE:

The security assessment was scoped to the following smart contract:

- Undelegation.sol
- Timelock.sol
- Staking.sol
- Rewards.sol
- Ownable.sol
- NodeProxy.sol

Audit Commit ID :

- eef82c8d5252f56d3357dd1ba4c1fc788e7faabd

Fixed Commit ID :

- b04fc3be788a6d698071ceb77f6fe844e0ded0e7

Fixed Updated Commit ID:

- c88a979fabb1f1338683d2687155558f7b006b4c

2. ASSESSMENT SUMMARY & FINDINGS OVERVIEW

CRITICAL	HIGH	MEDIUM	LOW	INFORMATIONAL
0	0	1	5	11

LIKELIHOOD

EXECUTIVE OVERVIEW

SECURITY ANALYSIS	RISK LEVEL	REMEDIATION DATE
HALØ1 – TIMELOCK CAN BE BYPASSED	Medium	SOLVED - 11/05/2022
HAL02 – LACK OF TRANSFEROWNERSHIP PATTERN	Low	SOLVED - 11/05/2022
HAL03 - DIFFERENT PROXYNODE ARRAY LENGTH REQUIREMENTS	Low	SOLVED - 11/04/2022
HAL04 – MISSING ZERO ADDRESS CHECKS	Low	SOLVED - 11/04/2022
HAL05 - LACK OF PARAMETER LIMITS	Low	PARTIALLY SOLVED - 11/04/2022
HALØ6 – MISSING REENTRANCY GUARD	Low	SOLVED - 11/04/2022
HAL07 - NODEPROXY ARRAY CANNOT BE MODIFIED	Informational	ACKNOWLEDGED
HAL08 - FLOATING PRAGMA	Informational	SOLVED - 11/04/2022
HAL09 - CACHE ARRAY LENGTH IN FOR LOOPS CAN SAVE GAS	Informational	SOLVED - 11/04/2022
HAL10 - USE CUSTOM ERRORS INSTEAD OF REVERT STRINGS	Informational	PARTIALLY SOLVED - 11/05/2022
HAL11 - REVERT STRING SIZE OPTIMIZATION	Informational	ACKNOWLEDGED
HAL12 – UNUSED EVENTS	Informational	SOLVED - 11/04/2022
HAL13 - UNNECESSARY CHECK	Informational	ACKNOWLEDGED
HAL14 – NO NEED TO INITIALIZE VARIABLES WITH DEFAULT VALUES	Informational	ACKNOWLEDGED
HAL15 - USING POSTFIX OPERATORS IN LOOPS	Informational	ACKNOWLEDGED
HAL16 - DIVISION BY ZERO	Informational	SOLVED - 11/04/2022
HAL17 - SPLITTING REQUIRE() STATEMENTS THAT USES AND OPERATOR SAVES GAS	Informational	SOLVED - 11/04/2022

FINDINGS & TECH DETAILS

3.1 (HAL-01) TIMELOCK CAN BE BYPASSED - MEDIUM

Description:

Timelock contract is used to queue the transfer of HBAR from the previous Staking contract to the new one. It introduces a lockedPeriod parameter, defining the minimum time between a withdrawal is requested and when it can be completed.

However, it has been detected that the address defined in timelockOwner can either queue funds that could be transferred once lockedPeriod has passed by and also can call setLockedPeriod(), which defines the value of lockedPeriod.

Since the same user who calls queuePartialFunds() or queueAllFunds() can set lockedPeriod by calling setLockedPeriod(), Timelock functionalities can be trivially bypassed by an ill-intentioned user with enough privileges, defeating the whole purpose of the contract.

Code Location:

List	ing	1:	Timelock.sol (Lines 64,84)
61		fun	ction queuePartialFunds(address payable to, uint256 amount)
62			external
63			checkZeroAddress(to)
64			
65			returns (uint256)
66		{	
67			if (amount > address(this).balance) revert("Amount exceeds
L,	bal	ance	e");
68			uint256 index = withdrawQueue.length;
69			Withdraw memory withdrawData = Withdraw({
70			timestamp: block.timestamp,
			lockedAmount: amount,
			to: to
			});
74			withdrawQueue.push(withdrawData);

```
75 emit Queued(index, amount);
76 return index;
77 }
78
79 /// @notice queue the transaction for withdrawal of the entire
79 /// @param to address of the account to transfer the tokens to
80 /// @param to address of the account to transfer the tokens to
81 function queueAllFunds(address payable to)
82 external
83 checkZeroAddress(to)
84 checkOwner
85 returns (uint256)
86 {
87 uint256 index = withdrawQueue.length;
88 Withdraw memory userTransaction = Withdraw({
89 timestamp: block.timestamp,
90 lockedAmount: address(this).balance,
91 to: to
92 });
93 withdrawQueue.push(userTransaction);
94 emit Queued(index, address(this).balance);
95 return index;
96 }
```

Listing 2: Timelock.sol (Line 146) 144 /// @notice Set the locking period for the transfer of tokens 145 /// @param _lockedPeriod time in secs for withholding transfer 145 transaction 146 function setLockedPeriod(uint256 _lockedPeriod) external 147 lockedPeriod = _lockedPeriod; 148 }

Risk Level:

Likelihood - 1 Impact - 5

Proof of Concept:

For this Proof of Concept, the user with Owner privileges in Timelock contract will extract immediately the whole balance of Staking contract. To do so, setLockedPeriod() will be called to set a 0 lock period, and then queueAllFunds() and withdraw() will be called consecutively, extracting the complete balance stored in the contract, defeating the purpose of a Time Lock.

```
>>> bypassTimeLocker()
```

Balance of Staking contract --> 1000000000000000000

```
Balance of user3 --> 100000000000000000000000
```

Current Locked period value --> 7200

Setting 0 Locked period...

Transaction sent: 0x8605a584c8304dca88f4b1f162a545680fe53f0b1705881e5d74ab2ba929736a Gas price: 0.0 gwei Gas limit: 600000000 Nonce: 51 Staking.setLockedPeriod confirmed Block: 15884358 Gas used: 13695 (0.00%)

Queuing transfer and immediatly calling withdraw...

```
contract_Staking.queueAllFunds(user3, {'from': owner}) -->
Transaction sent: 0x22199210a32f150c6ae82ab9e29575505b19977c7e6c8823701d3dec41218116
Gas price: 0.0 gwei Gas limit: 60000000 Nonce: 52
Staking.queueAllFunds confirmed Block: 15884359 Gas used: 106053 (0.02%)
```

```
contract_Staking.withdraw(0, {'from': owner}) -->
Transaction sent: 0x7dd4c71fb51aa219ce3c7aa1c1dc26cb85565e0d7be6f937dee5677e1b4a41ec
Gas price: 0.0 gwei Gas limit: 60000000 Nonce: 53
Staking.withdraw confirmed Block: 15884360 Gas used: 26791 (0.00%)
```

Balance of Staking contract --> 0

Balance of user3 --> 100001000000000000000000

Recommendation:

Multiple measures are recommended to enforce the correct use of Timelock:

- Separate the roles in charge of setting lockedPeriod and transferring balances.
- Protect setLockedPeriod() function itself with a Time Lock.
- Add a require statement containing the minimum acceptable value for lockedPeriod.

If this finding poses no security risk at all, then deleting Timelock contract is recommended for saving gas and using a regular transfer function instead.

Remediation Plan:

SOLVED: The Stader Labs team solved this finding by blocking the owner to set the lockedPeriod variable less than 2 days.

3.2 (HAL-02) LACK OF TRANSFEROWNERSHIP PATTERN - LOW

Description:

The current ownership transfer process for Timelock contract involves the current owner calling the setTimeLockOwner() function:

```
Listing 3: Timelock.sol

134 /// @notice Set new multisig owner for the transfer of Hbar to

L, new version

135 /// @param _timelockOwner the new owner of Hbar withdrawal to

L, new version

136 function setTimeLockOwner(address _timelockOwner)

137 external

138 checkZeroAddress(_timelockOwner)

139 checkOwner

140 {

141 timelockOwner = _timelockOwner;

142 }
```

If the nominated EOA account is not a valid account, it is entirely possible that the owner may accidentally transfer ownership to an uncontrolled account, losing the access to all functions with the checkOwner modifier.

Risk Level:

Likelihood - 2 Impact - 2

Recommendation:

It is recommended to implement a two-step process where the owner nominates an account and the nominated account needs to call an acceptOwnership() function for the transfer of the ownership to fully succeed. This ensures the nominated EOA account is a valid and active account.

Remediation Plan:

SOLVED: The Stader Labs team solved this finding by implementing a two-step process.

3.3 (HAL-03) DIFFERENT PROXYNODE ARRAY LENGTH REQUIREMENTS - LOW

Description:

It has been observed that two functions check the length of the input array and compare it to the length of nodeProxyAddresses array in Staking contract, but they do it differently.

collectRewards takes an input array which defines from which contracts rewards will be collected, and a require statement enforces that the length of the array is equal than nodeProxyAddresses length:

Lis	ting 4	4:	Staking.sol (Line 294)
286	f	unc	tion collectRewards(uint256[] memory
Ļ	pendi	ngR	ewardNodeIndexes)
287			external
288			payable
289			whenNotPaused
290			onlyOperator
291	{		
292			<pre>require(nodeStakingActive, "node staking not active");</pre>
293			require(
294			<pre>pendingRewardNodeIndexes.length == nodeProxyAddresses.</pre>
Ļ	lengt	h,	
295			"Invalid pendingRewardNodeIndexes input"
296);
297			<pre>for (uint256 i; i < nodeProxyAddresses.length; i++) {</pre>
298			<pre>if (pendingRewardNodeIndexes[i] == 1) {</pre>
299			require (
300			address(this).balance >= 1,
301			"Insufficient balance to execute
	colle	ctR	ewards"
302);
303			<pre>moveBalanceForStaking(nodeProxyAddresses[i], 1);</pre>
304			}
305			}
306	}		

On the other hand, stakeWithNodes takes as input a similar array in which each index of the position corresponds with each index of the nodeProxyAddresses array. However, this time, the require statement only checks for that array to have the same or lower length than nodeProxyAddresses:

Listing 5: Staking.sol (Line 247)

239	fun	ction stakeWithNodes(uint256[] calldata amountToSend,
L,	uint256	index)
240		external
241		whenNotPaused
242		onlyOperator
243	{	
244		<pre>require(!nodeStakingActive, "node staking already active")</pre>
Ļ	;	
245		<pre>require(index < nodeProxyAddresses.length, "Invalid index"</pre>
L,);	
246		require(
247		<pre>amountToSend.length <= nodeProxyAddresses.length,</pre>
248		"Invalid size of amountToSend"
249);
250		<pre>nodeStakingActive = true;</pre>
251		<pre>balanceBefore = address(this).balance;</pre>
252		<pre>// iterating over amountToSend array to send hbar to</pre>
	respect	ive index of nodeProxyContract
253		// following checks are to incorporate changes in the
	staking	contract balance after computing amountToSend
254		for (uint256 1 = 0; 1 < amountloSend.length; 1++) {
255		lt (
250		amountlosena[1] > 0 & &
257		address(this) balance $> 0 & &$
250		
260		/ L moveBalanceForStaking(nodeProvvAddresses[i]
	amountTo	oSend[i]):
261		} else if (
262		amountToSend[i] > 0 &&
263		address(this).balance > 0 &&
264		amountToSend[i] > address(this).balance
265) {
266		moveBalanceForStaking(
267		nodeProxyAddresses[i],
268		address(this).balance

```
269 );
270 }
271 }
272 if (address(this).balance > 0) {
273 moveBalanceForStaking(
274 nodeProxyAddresses[index],
275 address(this).balance
276 );
277 }
278 emit stakedWithNodes(balanceBefore);
279 }
280
```

Having two arrays in which each position corresponds with the same position of nodeProxyAddresses array but with different length requirements might be confusing and error-prone, since the absence of a single position in amountToSend could mean that every amount defined in the array is staked into the wrong node.

Risk Level:

Likelihood - 1 Impact - 3

Recommendation:

It is recommended to unify the criteria of the require statements on both functions, enforcing that amountToSend array has the same length of nodeProxyAddresses to prevent confusions or any input error.

Remediation Plan:

SOLVED: The Stader Labs team solved this finding by requiring amountToSend to have the same length as nodeProxyAddresses.

3.4 (HAL-04) MISSING ZERO ADDRESS CHECKS - LOW

Description:

The constructor of the Rewards.sol contract is missing address validation. Every address should be validated and checked that it is different from zero. Control of that constructor is wrong because of the or statement between them. Because of that issue, either the <u>stakerAddress</u> variable or <u>daoAddress</u> might be a 0 address, and it can cause an unintended loss in the <u>distributeStakingRewards</u> function. This is also considered a best practice.

Code Location:

Lis	ting 6: Rewards.sol (Line 61)
59 ∟ 60	<pre>constructor(address payable _stakerAddress, address payable _daoAddress) { require(</pre>
61	_stakerAddress != address(0) _daoAddress != address
Ļ	(0),
62	"Address cannot be a zero"
63);
64	stakerAddress = _stakerAddress;
65	daoAddress = _daoAddress;
66	genesisTimestamp = block.timestamp;
67	lastRedeemedTimestamp = genesisTimestamp;
68	<pre>// _pause();</pre>
69	}

Risk Level:

Likelihood - 2 Impact - 2

Recommendation:

It is recommended to validate that every address input is different from zero.

Remediation Plan:

SOLVED: The Stader Labs team solved this finding by enforcing separate zero address checks on each of the input addresses.

3.5 (HAL-05) LACK OF PARAMETER LIMITS - LOW

Description:

It has been detected that some parameter modifying functions do not have logical limits. This may cause the contract to function with parameter values that, although allowed, make no sense in the context of the application, which might cause a variety of problems or even rendering the contract unusable.

There are two functions in Undelegation.sol and Timelock.sol contracts that should have a minimum value check in place. These functions determine the minimum time needed for being able to unstake HBAR from Staking contract and for transferring funds from the old staking contract to the new one, respectively.

Having no minimum value check means that **HBAR** could be immediately unstaked or Timelock could be bypassed.

Similarly, it is recommended to define some boundaries on Staking.sol's minDeposit and maxDeposit, since setting a minDeposit value too high or a maxDeposit value too low (or zero) would prevent anyone from being able to stake HBAR.

Code Location:

```
Listing 7: "Timelock.sol

146 function setLockedPeriod(uint256 _lockedPeriod) external

L, checkOwner {

147 lockedPeriod = _lockedPeriod;

148 }
```

```
Listing 8: "Undelegation.sol
```

```
97 function setUnbondingTime(uint256 _unbondingTime) external

L onlyOwner {

98 unbondingTime = _unbondingTime;

99 emit NewUnbondingTime(unbondingTime);

100 }
```

Listing 9: "Staking.sol

```
393 /// @notice Set minimum deposit amount (onlyOwner)
394 /// @param _newMinDeposit the minimum deposit amount in
4 multiples of 10**8
395 function updateMinDeposit(uint256 _newMinDeposit) external
4 onlyOwner {
396 minDeposit = _newMinDeposit;
397 }
398
399 /// @notice Set maximum deposit amount (onlyOwner)
400 /// @param _newMaxDeposit the maximum deposit amount in
401 function updateMaxDeposit(uint256 _newMaxDeposit) external
402 maxDeposit = _newMaxDeposit;
403 }
```

Risk Level:

Likelihood - 2 Impact - 2

Recommendation:

It is recommended to enforce logical value limits for critical parameters and check for additional occurrences of this same vulnerability.

Remediation Plan:

PARTIALLY SOLVED: The Stader Labs team partially solved this finding by adding some logical checks on the Staking.sol contract, enforcing that minDeposit is lower than maxDeposit, and maxDeposit is greater than minDeposit.

3.6 (HAL-06) MISSING REENTRANCY GUARD - LOW

Description:

To protect against cross-function re-entrancy attacks, it may be necessary to use a mutex. By using this lock, an attacker can no longer exploit the withdrawal function with a recursive call. OpenZeppelin has its own mutex implementation called **ReentrancyGuard** which provides a modifier to any function called **nonReentrant** that guards the function with a mutex against re-entrancy attacks.

Code Location:

```
Listing 10: Timelock.sol
```

100	function withdraw(uint256 index) external returns (uint256) {
101	if (address(this).balance == 0) revert("No funds to
\vdash	withdraw");
102	<pre>if (index >= withdrawQueue.length) revert("Invalid index")</pre>
Ļ	;
103	Withdraw storage withdrawData = withdrawQueue[index];
104	if (withdrawData.timestamp + lockedPeriod >= block.
$ \downarrow $	timestamp)
105	<pre>revert("Unlock period not expired");</pre>
106	<pre>if (withdrawData.lockedAmount == 0) revert("Amount not</pre>
$ \vdash $	available");
107	address payable to = withdrawData.to;
108	<pre>uint256 amount = withdrawData.lockedAmount;</pre>
109	<pre>delete withdrawQueue[index];</pre>
110	<pre>// payable(to).transfer(amount);</pre>
	Address.sendValue(payable(to), amount);
	<pre>emit Transferred(index, amount, to);</pre>
	return index;
	}

Risk Level:

Likelihood - 2 Impact - 2

Recommendation:

The functions on the code location section have missing nonReentrant modifiers. It is recommended to add the OpenZeppelin ReentrancyGuard library to the project and use the nonReentrant modifier to avoid introducing future re-entrancy vulnerabilities.

Remediation Plan:

SOLVED: The Stader Labs team solved this finding by adding the nonReentrant modifier to the withdraw() function on the Timelock.sol contract.

3.7 (HAL-07) NODEPROXY ARRAY CANNOT BE MODIFIED - INFORMATIONAL

Description:

For each node in Hedera network, a NodeProxy contract will be deployed. Each one of these contracts is assigned to one node, and staking on one of the contracts will effectively mean the same as staking on the node itself.

The addresses of these contracts are stored in nodeProxyAddresses array, set in Staking contract's constructor function.

However, there is no way to modify the addresses contained in the array, so any eventuality related with Hedera nodes might potentially render Staking contract partially or even completely unusable. These eventualities might be related to node additions, node ID changes, etc.

Code Location:

Listin	g 11:	Staking.sol (Line 110)
104	con	structor(
105		address _hbarxAddress,
106		address _multisigAdminAddress,
107		address payable _undelegationContractAddress,
108		uint256 _totalSupply,
109		address _operator,
110		address[] memory _nodeProxyAddresses
111)	
112		Timelock(_multisigAdminAddress)
113		checkZeroAddress(_hbarxAddress)
114		<pre>checkZeroAddress(_undelegationContractAddress)</pre>
115		checkZeroAddress(_operator)
116	{	
117		hbarxAddress = _hbarxAddress;
118		<pre>undelegationContractAddress = _undelegationContractAddress</pre>
↓;		
119		<pre>totalSupply = _totalSupply;</pre>

```
120 operator = _operator;
121 for (uint256 i = 0; i < _nodeProxyAddresses.length; i++) {
122 if (_nodeProxyAddresses[i] == address(0))
123 revert("zero address for nodeProxy");
124 nodeProxyAddresses.push(payable(_nodeProxyAddresses[i
L, ]));
125 }
126 }
```

Risk Level:

Likelihood - 1 Impact - 2

Recommendation:

It is recommended to implement a protected function allowing Stader Labs team to modify nodeProxyAddresses array, adding flexibility in the eventuality of any node change on Hedera network.

Remediation Plan:

ACKNOWLEDGED: The Stader Labs team acknowledged this finding.

3.8 (HAL-08) FLOATING PRAGMA -INFORMATIONAL

Description:

Hedera Stader Protocol contracts use floating pragma. Contracts should be deployed with the same compiler version and flags they have tested thoroughly. Locking the pragma helps to ensure that contracts do not accidentally get deployed using, for example, either an outdated compiler version that might introduce bugs that affect the contract system negatively or a pragma version that is too new and has not been extensively tested.

Risk Level:

Likelihood - 1 Impact - 1

Recommendation:

Consider locking the pragma version with known bugs for the compiler version. When possible, do not use floating pragma in the final live deployment. Specifying a fixed compiler version ensures that the bytecode produced does not vary between builds. This is especially important if you rely on bytecode-level verification of the code.

Remediation Plan:

SOLVED: The Stader Labs team solved this finding by fixing the pragma version for each contract in scope.

3.9 (HAL-09) CACHE ARRAY LENGTH IN FOR LOOPS CAN SAVE GAS -INFORMATIONAL

Description:

Reading array length at each iteration of the loop takes 6 gas (3 for mload and 3 to place memory_offset) in the stack. Caching the array length in the stack saves around 3 gas per iteration.

Code Location:

Lis	ting 12:	Staking.sol (Lines 294,297)
286	fund	ction collectRewards(uint256[] memory
287	penuingi	external
288		payable
289		whenNotPaused
290		onlyOperator
291	{	
292		<pre>require(nodeStakingActive, "node staking not active");</pre>
293		require(
294		<pre>pendingRewardNodeIndexes.length == nodeProxyAddresses.</pre>
4	length,	
295		"Invalid pendingRewardNodeIndexes input"
296);
297		<pre>for (uint256 i; i < nodeProxyAddresses.length; i++) {</pre>

Listing 13: Staking.sol (Line 314)

```
312 function withdrawFromNodes() external whenNotPaused

   onlyOperator {
313 require(nodeStakingActive, "node staking not active");
314 for (uint256 i; i < nodeProxyAddresses.length; i++) {</pre>
```

Risk Level:

Likelihood - 1 Impact - 1

Recommendation:

Consider the cache array length.

Remediation Plan:

SOLVED: The Stader Labs team solved this finding by caching the array length when needed.

3.10 (HAL-10) USE CUSTOM ERRORS INSTEAD OF REVERT STRINGS -INFORMATIONAL

Description:

Starting from Solidity v0.8.4, there is a convenient and gas-efficient way to explain to users why an operation failed through the use of custom errors. If the revert string uses strings to provide additional information about failures (e.g. require(!isStakePaused, 'Staking is paused');), but they are rather expensive, especially when it comes to deploying cost, and it is difficult to use dynamic information in them.

Risk Level:

Likelihood - 1 Impact - 1

Recommendation:

It is recommended to implement custom errors instead of revert strings.

Remediation Plan:

PARTIALLY SOLVED: The Stader Labs team partially solved this finding by changing revert strings to custom errors in a few files.

3.11 (HAL-11) REVERT STRING SIZE OPTIMIZATION - INFORMATIONAL

Description:

Shortening the revert strings to fit within 32 bytes will decrease deployment time gas and decrease runtime gas when the revert condition is met.

Revert strings that are longer than 32 bytes require at least one additional mstore, along with additional overhead to calculate memory offset, etc. For example:

Code Location:

Listing	14: Rewards.sol
82	require(
83	address(this).balance > 0,
84	"Contract balance is should be greater than 0"
85);

Risk Level:

Likelihood - 1 Impact - 1

Recommendation:

Shorten the revert strings to fit within 32 bytes. That will affect gas optimization.

Remediation Plan:

ACKNOWLEDGED: The Stader Labs team acknowledged this finding.

3.12 (HAL-12) UNUSED EVENTS -INFORMATIONAL

Description:

The following events are declared, but they are not emitted by any function:

Code Location:

NodeProxy.sol
- Line 17:
event Received(address indexed from, uint256 amount);
- Line 20:
event Fallback(address indexed from, uint256 amount);
- Line 29:
event CollectedRewards();

Staking.sol
- Line 70:
event Undelegated(address indexed to, uint256 amount);

Undelegation.sol
- Line 27:
event Received(address from, uint256 amount);
- Line 29:
event Fallback(address from, uint256 amount);

Risk Level:

Likelihood - 1 Impact - 1

Recommendation:

Check whether these events should be used and if not remove them.

Remediation Plan:

SOLVED: The Stader Labs team solved this finding by removing unnecessary events.

3.13 (HAL-13) UNNECESSARY CHECK - INFORMATIONAL

Description:

The distributeStakingRewards function under Rewards.sol does check if the daoFeesPercentage is less than 100 and then reverts. However, that condition will never be reachable as the setDaoFeesPercentage function guarantees that daoFeesPercentage won't be able to higher or equal to 100.

Code Location:

List	ting 15: Rewards.sol (Lines 86,87,88,89)
81	<pre>function distributeStakingRewards() external whenNotPaused</pre>
Ļ	nonReentrant {
82	require(
83	address(this).balance > 0,
84	"Contract balance is should be greater than 0"
85);
86	require(
87	daoFeesPercentage < 100,
88	"Dao fees percentage should be less than 100"
89);

Listing	16: Rewards.sol (Line 151)
146	<pre>function setDaoFeesPercentage(uint256 _daoFeesPercentage)</pre>
147	external
148	onlyOwner
149	{
150	require(
151	_daoFeesPercentage < 100,
152	"Dao fees percentage should be less than 100"
153);
154	daoFeesPercentage = _daoFeesPercentage;
155	}

Risk Level:

Likelihood - 1 Impact - 1

Recommendation:

It is recommended to remove unnecessary checks to reduce gas costs

Remediation Plan:

ACKNOWLEDGED: The Stader Labs team acknowledged this finding.

3.14 (HAL-14) NO NEED TO INITIALIZE VARIABLES WITH DEFAULT VALUES -INFORMATIONAL

Description:

uint256 variables are already initialized to 0 by default. uint256 public epoch = 0 would reassign the 0 to epoch which wastes gas.

The same occurs with bool and address variables. They are already initialized to false/address(0).

Code Location:

- Line 14:

```
Rewards.sol
- Line 21:
uint256 public epoch = 0;
```

```
Staking.sol
- Line 30:
bool public isStakePaused = false;
- Line 31:
bool public isUnstakePaused = false;
- Line 32:
bool public nodeStakingActive = false;
- Line 36:
uint256 public minDeposit = 0;
- Line 39:
uint256 public totalSupply = 0;
- Line 121:
for (uint256 i = 0; i < _nodeProxyAddresses.length; i++){</pre>
- Line 254:
for (uint256 i = 0; i < amountToSend.length; i++){</pre>
NodeProxy.sol
```

address payable public stakerAddress = payable(address(0));

Risk Level:

Likelihood - 1 Impact - 1

Recommendation:

It is recommended to not initialize uint variables to 0. bool variables to false and address variables to address(0) to save some gas. For example, use instead: uint256 public totalSupply;

Remediation Plan:

ACKNOWLEDGED: The Stader Labs team acknowledged this finding.

3.15 (HAL-15) USING POSTFIX OPERATORS IN LOOPS - INFORMATIONAL

Description:

In the loops below, postfix operators (e.g. i++) were used to increment or decrement the value of variables. In loops, using prefix operators (e.g., ++i) costs less gas per iteration than postfix operators.

Code Location:

```
Staking.sol
- Line 121:
for (uint256 i = 0; i < _nodeProxyAddresses.length; i++){
- Line 254:
for (uint256 i = 0; i < amountToSend.length; i++){
- Line 297:
for (uint256 i; i < nodeProxyAddresses.length; i++){
- Line 314:
for (uint256 i; i < nodeProxyAddresses.length; i++){</pre>
```

Risk Level:

Likelihood - 1 Impact - 1

Recommendation:

It is recommended to use ++i instead of i++ to increment the value of an uint variable inside a loop. This does not only apply to the iterator variable. It also applies to the increments/decrements done inside the loop code block.

Remediation Plan:

ACKNOWLEDGED: The Stader Labs team acknowledged this finding.

3.16 (HAL-16) DIVISION BY ZERO -INFORMATIONAL

Description:

Calling the getExchangeRate function with totalSupply as 0 and nodeStakingActive as true, will cause the function to throw a division by zero error.

Code Location:

```
Listing 17: "Staking.sol (Lines 367,368)

function getExchangeRate() external view returns (uint256) {

///@dev 1HBar = 100_000_000 tinybar

if (nodeStakingActive) {

return (balanceBefore * decimals) / totalSupply;

}

uint256 exchangeRate = 1 * decimals;

if (totalSupply == 0 || address(this).balance == 0) {

return exchangeRate;

} else {

exchangeRate = ((address(this).balance) * decimals) /

totalSupply;

}

constant exchangeRate;

}

f return exchangeRate;

}

f return exchangeRate = ((address(this).balance) * decimals) /

totalSupply;

}

}
```

Risk Level:

Likelihood - 1 Impact - 1

Recommendation:

Make sure to validate all operands used during a math operation and inform the user of unappropriated state by reverting the transaction with

a custom message.

Remediation Plan:

SOLVED: The Stader Labs team solved this finding by validating all operands before performing math operations.

3.17 (HAL-17) SPLITTING REQUIRE() STATEMENTS THAT USES AND OPERATOR SAVES GAS - INFORMATIONAL

Description:

Instead of using the && operator in a single require statement to check multiple conditions, using multiple require statements with one condition per require statement will save 8 GAS per &&

The gas difference would only be realized if the revert condition is realized (met).

Code Location:

Lis	ting 18: Staking.sol
138	require(
139	hbarReceived > minDeposit && hbarReceived <=
Ļ	maxDeposit,
140	"Deposit amount must be within valid range"
141);

Listing 19: Undelegation.sol

62		require(
63		undelegateData.amount != 0 && undelegateData.timestamp
L,	!= 0,	
64		"Undelegation not found"
65);

Proof of Concept:

The following tests were carried out in remix with both optimization turned on and off

```
Listing 20
```

```
1 require ( a > 1 && a < 5, "Initialized");
2 return a + 2;</pre>
```

Execution cost 21617 with optimization and using &&

21976 without optimization and using &&

After splitting the require statement

Listing 21

```
1 require (a > 1 ,"Initialized");
2 require (a < 5 , "Initialized");
3 return a + 2;
```

Execution cost 21609 with optimization and split require 21968 without optimization and using split require

Risk Level:

Likelihood - 1 Impact - 1

Recommendation:

For best security practices, consider as much as possible, declaring events at the end of the function. Events can be used to detect the end of the operation.

Remediation Plan:

SOLVED: The Stader Labs team solved this finding by splitting require statements.

MANUAL TESTING

4.1 INTRODUCTION

Halborn performed different manual tests in all the different Facets of the **Hedera protocol**, trying to find any logic flaws and vulnerabilities.

During the manual tests, the following areas were reviewed carefully :

- 1. Reward mechanism.
- 2. Withdraw funds mechanism.
- 3. Undelegation mechanism.

4.2 TESTING

REWARD MECHANISM:

In the Rewards.sol contract there is a function named distributeStakingRewards and this function is responsible to distribute staking rewards among those with staker address and DAO address.

```
Listing 22: Rewards.sol
       function distributeStakingRewards() external whenNotPaused
           require(
               address(this).balance > 0,
           );
           require(
               daoFeesPercentage < 100,</pre>
           );
           uint256 currentTimestamp = block.timestamp;
           uint256 epochDelta = (currentTimestamp -
↓ lastRedeemedTimestamp);
           lastRedeemedTimestamp = currentTimestamp;
           epoch++;
           uint256 epochRewards = (epochDelta * emissionRate);
           uint256 totalRewards = address(this).balance;
           if (epochRewards > totalRewards) epochRewards =

    totalRewards; // this is important

           uint256 daoFees = (epochRewards * daoFeesPercentage) /
L→ 100;
           Address.sendValue(payable(stakerAddress), epochRewards -
emit DistributedRewards(
               stakerAddress,
               epochRewards - daoFees,
           );
```

```
108  // payable(daoAddress).transfer(daoFees);
109  Address.sendValue(payable(daoAddress), daoFees);
110  emit DaoTransfer(daoAddress, daoFees, currentTimestamp);
111 }
```

This function does not have any kind of msg.sender control, So anyone would be able to call this function. This function distributes rewards every 24 hours. A malicious actor can call this function before 24 hours, but a malicious actor can call this anytime. However, he will not be able to manipulate this function somehow because of the correctness of totalRewards and daoFees calculations. Even malicious actor calls this in 12 hours, rewards will be the same because it always calculates it with lastRedemeedTiestamp variable.

```
>>> lastRedeemedTimestamp = rewardsContract.lastRedeemedTimestamp()
>>> stakingContract.balance()
>>> rewardsContract.balance()
>>> tx1 = rewardsContract_distributeStakingRewards({'from': user1})
Transaction sent: 0x5cf4ea76175d8a4e62494e1afc08d4fc68341ba096b587937e5f0b4d13b42bd6
 Gas price: 0.0 gwei Gas limit: 12000000 Nonce: 0
 Rewards.distributeStakingRewards confirmed Block: 15 Gas used: 77636 (0.65%)
>>> stakingContract.balance()
83516507676
>>> epochDelta = tx1.timestamp - lastRedeemedTimestamp
>>> secondLastRedeemedTimestamp = rewardsContract.lastRedeemedTimestamp()
>>> tx2 = rewardsContract.distributeStakingRewards({'from': user1})
Transaction sent: 0xc49a2ee762d3ed47da0535b0df8c3728219b6b375032112e82d1971b9cf8c6aa
 Gas price: 0.0 gwei Gas limit: 12000000 Nonce: 1
 Rewards.distributeStakingRewards confirmed Block: 16
                                                        Gas used: 62636 (0.52%)
>>> stakingContract.balance()
272820591742
 >> ongchDaltoChauldDa = tx2 timestame_plastBadeemodTimestame
>>> epochDeltaShouldBe = tx2.timestamp - lastRedeemedTimestamp
>>> secondEpochDelta = tx2.timestamp - secondLastRedeemedTimestamp
>>> epochDeltaShouldBe == secondEpochDelta + epochDelta
True
>>> epochRewards = epochDeltaShouldBe * rewardsContract.emissionRate()
>>> if epochRewards > initialRewardsBalance:
          .
epochRewards = initialRewardsBalance
>>> daoFees = epochRewards * rewardsContract.daoFeesPercentage() / 100
>>> stakerShouldGet = epochRewards - daoFees
```

Listing 23

```
1 >>> stakerShouldGet
2 272820591741.6
```

```
3 >>> stakingContract.balance()
```

4 272820591742

So even attacker calls this function 2 times, in the end staker gets the same balance because the math is correct. The only difference here is the precision loss in division operation in solidity. The only way to exploit this function is to make currentTimestamp variable to be equal to lastRedeemedTimestamp and if an attacker can do that. epochDelta will be 0 and eventually staker or DAO won't be able to get any rewards. But it's not likely possible because the attacker can't send thousands of transactions to make it pass in the same block time. The average block time mining in Ethereum is 12 seconds, so the attacker won't be able to do this.

WITHDRAW FUNDS MECHANISM:

In the Timelock.sol contract there is a function named withdraw, with using this function users can withdraw their funds. And there are two other functions named queuePartialFunds and queueAllFunds which allow the owner to queue funds.

Listing 24: Timelock.sol

61	function queuePartialFunds(address payable to, uint256 amount)
62	external
63	checkZeroAddress(to)
64	checkOwner
65	returns (uint256)
66	{
67	if (amount > address(this).balance) revert("Amount exceeds
Ļ	balance");
68	uint256 index = withdrawQueue.length;
69	Withdraw memory withdrawData = Withdraw({
70	timestamp: block.timestamp,
	lockedAmount: amount,
	to: to
	});
74	withdrawQueue. push(withdrawData);
75	<pre>emit Queued(index, amount);</pre>
76	return index;
	}

Listing 25: Timelock.sol

81	<pre>function queueAllFunds(address payable to)</pre>
82	external
83	checkZeroAddress(to)
84	checkOwner
85	returns (uint256)
86	{
87	uint256 index = withdrawQueue.length;
88	Withdraw memory userTransaction = Withdraw({
89	timestamp: block.timestamp,
90	<pre>lockedAmount: address(this).balance,</pre>
91	to: to
92	});
93	withdrawQueue.push(userTransaction);

```
94 emit Queued(index, address(this).balance);
95 return index;
96 }
```

```
Listing 26: Timelock.sol, (Lines 109,111)
       function withdraw(uint256 index) external returns (uint256) {
           if (address(this).balance == 0) revert("No funds to
 ↓ withdraw");
           if (index >= withdrawQueue.length) revert("Invalid index")
           Withdraw storage withdrawData = withdrawQueue[index];
           if (withdrawData.timestamp + lockedPeriod >= block.
 \downarrow  timestamp)
               revert("Unlock period not expired");
           if (withdrawData.lockedAmount == 0) revert("Amount not
address payable to = withdrawData.to;
           uint256 amount = withdrawData.lockedAmount;
           delete withdrawQueue[index];
           Address.sendValue(payable(to), amount);
           emit Transferred(index, amount, to);
```

At the end of the withdraw function, contract is sending withdrawData. lockedAmount amount to withdrawData.to user. Address.sendValue is sending the given Ethereum amount to the user with .call function.

```
Listing 27: Address.sol, (Line 63)
60 function sendValue(address payable recipient, uint256 amount)
L, internal {
61 require(address(this).balance >= amount, "Address:
L, insufficient balance");
62
63 (bool success, ) = recipient.call{value: amount}("");
64 require(success, "Address: unable to send value, recipient
L, may have reverted");
65 }
```

So in theory it is possible to create another contract to make a re-entrancy attack because withdraw function does not have any re-entrancy guard mechanism. However, withdraw function follows the Check-Effects -Interaction pattern correctly because it is deleting the index before sending the Ethereum to the user. So, it is not likely possible to do re-entracy attack in this function. There is no msg.sender control in this withdraw function, so that means any user can call this withdraw function with any parameter and withdraw funds for someone else's Ethereum to their address, but it is not a vulnerability because at the end the correct user is funded.

ומווכ נוכמו בא ווטו טכובווכט >>> stakingContract.queuePartialFunds(user1, web3.toWei('0.1', 'ether'), {'from': deployer})
Transaction sent: 0x22a314a69c0fdc126343f922593b70e85b9d44a09372908d23d4b6799ec56a49 Gas price: 0.0 gwei Gas limit: 12000000 Nonce: 16 Staking.queuePartialFunds confirmed Block: 19 Gas used: 106354 (0.89%) <Transaction '0x22a314a69c0fdc126343f922593b70e85b9d44a09372908d23d4b6799ec56a49'> >>> stakingContract.queuePartialFunds(user2, web3.toWei('0.1', 'ether'), {'from': deployer})
Transaction sent: 0x0c752e8b96cb0eacc7a63bc0f66773fa695b7849efd375695a3e5dfcb879b009 Gas price: 0.0 gwei Gas limit: 12000000 Nonce: 17 Staking.queuePartialFunds confirmed Block: 20 Gas used: 91342 (0.76%) <Transaction '0x0c752e8b96cb0eacc7a63bc0f66773fa695b7849efd375695a3e5dfcb879b009'> >>> stakingContract.withdrawQueue(0) (1667128388, 10000000000000000, "0x33A4622B82D4c04a53e170c638B944ce27cffce3") >>> stakingContract withdrawQueue(1) (1667128391, 1000000000000000000, "0x0063046686E46Dc6F15918b61AE2B121458534a5") >>> stakingContract.lockedPeriod() >>> stakingContract.setLockedPeriod(0, {'from': deployer}) Transaction sent: 0x9cfd78d1b30f7f617bb26e14c76638916def940d94eed6f00d2ece6d2ddbdf63 Gas price: 0.0 gwei Gas limit: 12000000 Nonce: 18 Staking.setLockedPeriod confirmed Block: 21 Gas used: 13695 (0.11%) <Transaction '0x9cfd78d1b30f7f617bb26e14c76638916def940d94eed6f00d2ece6d2ddbdf63'> >>> stakingContract.lockedPeriod()

```
>>> user1.balance()
1000000000000000000000000
>>> user2.balance()
1000000000000000000000000
>>> stakingContract.withdraw(0, {'from': user1})
Transaction sent: 0x654fb22e95687c073d770f273ee86caf486c940bd5a17e8749d4cc7d8531bbb5
Gas price: 0.0 gwei Gas limit: 12000000 Nonce: 2
Staking.withdraw confirmed Block: 22 Gas used: 26791 (0.22%)
```

UNDELEGATION MECHANISM:

In the Staking.sol contract after users unstake their **HBARX** to withdraw their money, unStake function calls the undelegate function of Undelegation.sol contract.

```
Listing 28: Staking.sol, (Lines 222,223,224)
189 function unStake(uint256 amount) external whenNotPaused returns (
→ uint256) {
           require(!nodeStakingActive, "node Staking is active");
           require(!isUnstakePaused, "Unstaking is paused");
           uint256 hbarxToBurn = (amount);
           uint256 hbarToTransfer = hbarxToBurn; // exchange rate = 1
           if (totalSupply != 0) {
                    (hbarxToBurn * ((address(this).balance))) /
                    (totalSupply);
           }
           int256 transferTokenResponse = HederaTokenService.
 ↓ transferToken(
               hbarxAddress,
               msg.sender,
               address(this),
               hbarxToBurn.toInt256().toInt64()
           );
           if (transferTokenResponse != HederaResponseCodes.SUCCESS)
_ ↓ {
                revert("HBARX transfer failed");
           (int256 burnTokenResponse, uint64 newTotalSupply) =
                .burnToken(hbarxAddress, hbarxToBurn.toUint64(), new
 \downarrow int64[](0));
           totalSupply = uint256(newTotalSupply);
           if (burnTokenResponse != HederaResponseCodes.SUCCESS) {
                revert("HBARX burn failed");
           }
```

```
220
221 ///@dev move tokens to undelegation contract
222 (bool success,) = payable(undelegationContractAddress).
L, call{
223 value: hbarToTransfer
224 }(abi.encodeWithSignature("undelegate(address)", msg.
L, sender));
225 if (!success) {
226 revert("Transfer failed");
227 }
228 emit UnStaked(msg.sender, hbarToTransfer, hbarxToBurn);
229 ///@dev return hbars for transaction
230 return hbarToTransfer;
231 }
```

The only way to call undelegate function is by calling the unStake function because undelegate function is checking if msg.sender is stakingContractAddress.

```
Listing 29: Undelegation.sol, (Lines 44,45,46,47,49)
```

```
42 function undelegate(address to) external payable returns (

43 require(msg.value > 0, "Undelegate amount must be greater

44 require(

45 msg.sender == stakingContractAddress,

46 "Only staking contract can undelegate"

47 );

48

49 undelegationsMap[to].push(Undelegate(block.timestamp, msg.

48

49 emit Undelegated(to, msg.value);

50 emit Undelegated(to, msg.value);

51 return msg.value;

52 }
```

And this function pushes given msg.value and to parameters to undelegationsMap for each user. After that step, users can call withdraw function to withdraw their money. withdraw function is calling Address.sendValue like in Timelock contract. However, in this case, there is a nonReentrant modifier to block the user to make re-entrancy attacks. Even without nonReentrant guard also on the function Check-Effects-Interaction pattern is used correctly.

```
Listing 30: Undelegation.sol, (Lines 60,72,74)
```

```
60 function withdraw(uint256 index) external whenNotPaused

nonReentrant {

61 Undelegate storage undelegateData = undelegationsMap[msg.

62 require(

63 undelegateData.amount != 0 && undelegateData.timestamp

64 "Undelegation not found"

65 );

66 require(

67 undelegateData.timestamp + unbondingTime <= block.

68 "Release time not reached"

69 );

70 

71 uint256 amount = undelegateData.amount;

72 delete undelegationsMap[msg.sender][index];

73 // payable(msg.sender).transfer(amount);

74 Address.sendValue(payable(msg.sender), amount);

75 emit Withdrawn(msg.sender, amount);

76 }

71
```

So, it is not possible to make re-entrancy attacks or get more money than you also deserve in this function.

```
>>> undelegationContract undelegate(user2, {'from': stakingContract, 'value': web3 toWei('0.1', 'ether')})
Transaction sent: 0x6e47dd21c9bb4eba35cf1666e54bfae50030ed3b10de14ba1455e2837caf276
Gas price: 0.0 gwei Gas Limit: 1200000 Nonce: 2
Undelegation.undelegate confirmed Block: 27 Gas used: 85184 (0.71%)

Transaction '0x6e47dd21c9bb4eba365cf1666e54bfae50030ed3b10de14ba1455e2837caf276'>
>>> undelegationContract.setUnbondingTime(0, {'from': deployer})
Transaction sent: 0x6771742e9bba3beae7796b28fb57ef849a097e3f9678bc5b56ce1d23581c09c4
Gas price: 0.0 gwei Gas Limit: 1200000 Nonce: 20
UndelegationContract.undelegate1onsMap(user2, 0)
(166713269, 100000000000000)
>>> undelegationContract.unbondingTime()
0
>>> undelegationContract.withdraw(0, {'from': user2})
Transaction sent: 0x7d7b6a47bab9a7f7c72910dc7aa56f26e4ebd85102889c9a6edd1b2466c86a94
Gas price: 0.0 gwei Gas Limit: 12000000 Nonce: 0
Undelegation.withdraw confirmed Block: 29 Gas used: 27053 (0.23%)
```

AUTOMATED TESTING

5.1 STATIC ANALYSIS REPORT

Description:

Halborn used automated testing techniques to enhance the coverage of certain areas of the smart contracts in scope. Among the tools used was Slither, a Solidity static analysis framework. After Halborn verified the smart contracts in the repositories and was able to compile them correctly into their ABIs and binary format, Slither was run against the contracts. This tool can statically verify mathematical relationships between Solidity variables to detect invalid or inconsistent usage of the contracts' APIs across the entire code-base.

Slither results:





• No major issues were found by Slither.



THANK YOU FOR CHOOSING